

Lecture 10

Basic Dividers

Required Reading

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Computer Arithmetic: Algorithms and Hardware Design

Chapter 13, Basic Division Schemes

13.1, Shift/Subtract Division Algorithms

13.3, Restoring Hardware Dividers

13.4, Non-Restoring and Signed Division

Chapter 15 Variation in Dividers

15.6, Combined Multiply/Divide Units

**Notation
and
Basic Equations**

Notation

z	Dividend	$z_{2k-1} z_{2k-2} \cdot \cdot \cdot z_2 z_1 z_0$
d	Divisor	$d_{k-1} d_{k-2} \cdot \cdot \cdot d_1 d_0$
q	Quotient	$q_{k-1} q_{k-2} \cdot \cdot \cdot q_1 q_0$
s	Remainder ($s = z - dq$)	$s_{k-1} s_{k-2} \cdot \cdot \cdot s_1 s_0$

Basic Equations of Division

$$z = d q + s$$

$$|s| < |d|$$

$$\text{sign}(s) = \text{sign}(z)$$

$$\begin{array}{l} z > 0 \\ 0 \leq s < |d| \end{array}$$

$$\begin{array}{l} z < 0 \\ -|d| < s \leq 0 \end{array}$$

Unsigned Integer Division Overflow

- Must check overflow because obviously the quotient q can also be $2k$ bits.
 - For example, if the divisor d is 1, then the quotient q is the dividend z , which is $2k$ bits

Condition for no overflow (i.e. q fits in k bits):

$$z = q d + s < (2^k - 1) d + d = d 2^k$$

$$z = z_H 2^k + z_L < d 2^k$$

$$z_H < d$$

Sequential Integer Division

Basic Equations

$$s^{(0)} = z$$

$$s^{(j)} = 2 s^{(j-1)} - q_{k-j} (2^k d) \quad \text{for } j=1..k$$

$$s^{(k)} = 2^k s$$

Sequential Integer Division Justification

$$s^{(1)} = 2 z - q_{k-1} (2^k d)$$

$$s^{(2)} = 2(2 z - q_{k-1} (2^k d)) - q_{k-2} (2^k d)$$

$$s^{(3)} = 2(2(2 z - q_{k-1} (2^k d)) - q_{k-2} (2^k d)) - q_{k-3} (2^k d)$$

.

$$\begin{aligned} s^{(k)} &= 2(\dots 2(2 z - q_{k-1} (2^k d)) - q_{k-2} (2^k d)) - q_{k-3} (2^k d) \dots \\ &\quad - q_0 (2^k d) = \\ &= 2^k z - (2^k d) (q_{k-1} 2^{k-1} + q_{k-2} 2^{k-2} + q_{k-3} 2^{k-3} + \dots + q_0 2^0) = \\ &= 2^k z - (2^k d) q = 2^k (z - d q) = 2^k s \end{aligned}$$

Fig. 13.2 Examples of sequential division with integer and fractional operands.

Integer division		Fractional division	
z	0 1 1 1 0 1 0 1	z_{frac}	. 0 1 1 1 0 1 0 1
2^4d	1 0 1 0	d_{frac}	. 1 0 1 0
$s^{(0)}$	0 1 1 1 0 1 0 1	$s^{(0)}$. 0 1 1 1 0 1 0 1
$2s^{(0)}$	0 1 1 1 0 1 0 1	$2s^{(0)}$	0 . 1 1 1 0 1 0 1
$-q_3 2^4d$	1 0 1 0 $\{q_3 = 1\}$	$-q_{-1}d$. 1 0 1 0 $\{q_{-1} = 1\}$
$s^{(1)}$	0 1 0 0 1 0 1	$s^{(1)}$. 0 1 0 0 1 0 1
$2s^{(1)}$	0 1 0 0 1 0 1	$2s^{(1)}$	0 . 1 0 0 1 0 1
$-q_2 2^4d$	0 0 0 0 $\{q_2 = 0\}$	$-q_{-2}d$. 0 0 0 0 $\{q_{-2} = 0\}$
$s^{(2)}$	1 0 0 1 0 1	$s^{(2)}$. 1 0 0 1 0 1
$2s^{(2)}$	1 0 0 1 0 1	$2s^{(2)}$	1 . 0 0 1 0 1
$-q_1 2^4d$	1 0 1 0 $\{q_1 = 1\}$	$-q_{-3}d$. 1 0 1 0 $\{q_{-3} = 1\}$
$s^{(3)}$	1 0 0 0 1	$s^{(3)}$. 1 0 0 0 1
$2s^{(3)}$	1 0 0 0 1	$2s^{(3)}$	1 . 0 0 0 1
$-q_0 2^4d$	1 0 1 0 $\{q_0 = 1\}$	$-q_{-4}d$. 1 0 1 0 $\{q_{-4} = 1\}$
$s^{(4)}$	0 1 1 1	$s^{(4)}$. 0 1 1 1
s	0 1 1 1	s_{frac}	0 . 0 0 0 0 0 1 1 1
q	1 0 1 1	q_{frac}	. 1 0 1 1

Fractional Division

Unsigned Fractional Division

z_{frac}	Dividend	$.z_{-1}z_{-2} \dots z_{-(2k-1)}z_{-2k}$
d_{frac}	Divisor	$.d_{-1}d_{-2} \dots d_{-(k-1)}d_{-k}$
q_{frac}	Quotient	$.q_{-1}q_{-2} \dots q_{-(k-1)}q_{-k}$
s_{frac}	Remainder	$\underbrace{.000\dots 0}_{k \text{ bits}}s_{-(k+1)} \dots s_{-(2k-1)}s_{-2k}$

Integer vs. Fractional Division

For Integers:

$$z = q d + s \quad | \cdot 2^{-2k}$$

$$z 2^{-2k} = (q 2^{-k}) (d 2^{-k}) + s (2^{-2k})$$

For Fractions:

$$z_{\text{frac}} = q_{\text{frac}} d_{\text{frac}} + s_{\text{frac}}$$

where

$$z_{\text{frac}} = z 2^{-2k}$$

$$d_{\text{frac}} = d 2^{-k}$$

$$q_{\text{frac}} = q 2^{-k}$$

$$s_{\text{frac}} = s 2^{-2k}$$

Unsigned Fractional Division Overflow

Condition for no overflow:

$$z_{\text{frac}} < d_{\text{frac}}$$

Sequential Fractional Division

Basic Equations

$$s^{(0)} = z_{\text{frac}}$$

$$s^{(j)} = 2 s^{(j-1)} - q_{-j} d_{\text{frac}} \quad \text{for } j=1..k$$

$$2^k \cdot s_{\text{frac}} = s^{(k)}$$

$$s_{\text{frac}} = 2^{-k} \cdot s^{(k)}$$

Sequential Fractional Division Justification

$$s^{(1)} = 2 z_{\text{frac}} - q_{-1} d_{\text{frac}}$$

$$s^{(2)} = 2(2 z_{\text{frac}} - q_{-1} d_{\text{frac}}) - q_{-2} d_{\text{frac}}$$

$$s^{(3)} = 2(2(2 z_{\text{frac}} - q_{-1} d_{\text{frac}}) - q_{-2} d_{\text{frac}}) - q_{-3} d_{\text{frac}}$$

.....

$$s^{(k)} = 2(\dots 2(2 z_{\text{frac}} - q_{-1} d_{\text{frac}}) - q_{-2} d_{\text{frac}}) - q_{-3} d_{\text{frac}} \dots$$

$$- q_{-k} d_{\text{frac}} =$$

$$= 2^k z_{\text{frac}} - d_{\text{frac}} (q_{-1} 2^{k-1} + q_{-2} 2^{k-2} + q_{-3} 2^{k-3} + \dots + q_{-k} 2^0) =$$

$$= 2^k z_{\text{frac}} - d_{\text{frac}} 2^k (q_{-1} 2^{-1} + q_{-2} 2^{-2} + q_{-3} 2^{-3} + \dots + q_{-k} 2^{-k}) =$$

$$= 2^k z_{\text{frac}} - (2^k d_{\text{frac}}) q_{\text{frac}} = 2^k (z_{\text{frac}} - d_{\text{frac}} q_{\text{frac}}) = 2^k s_{\text{frac}5}$$

Restoring Unsigned Integer Division

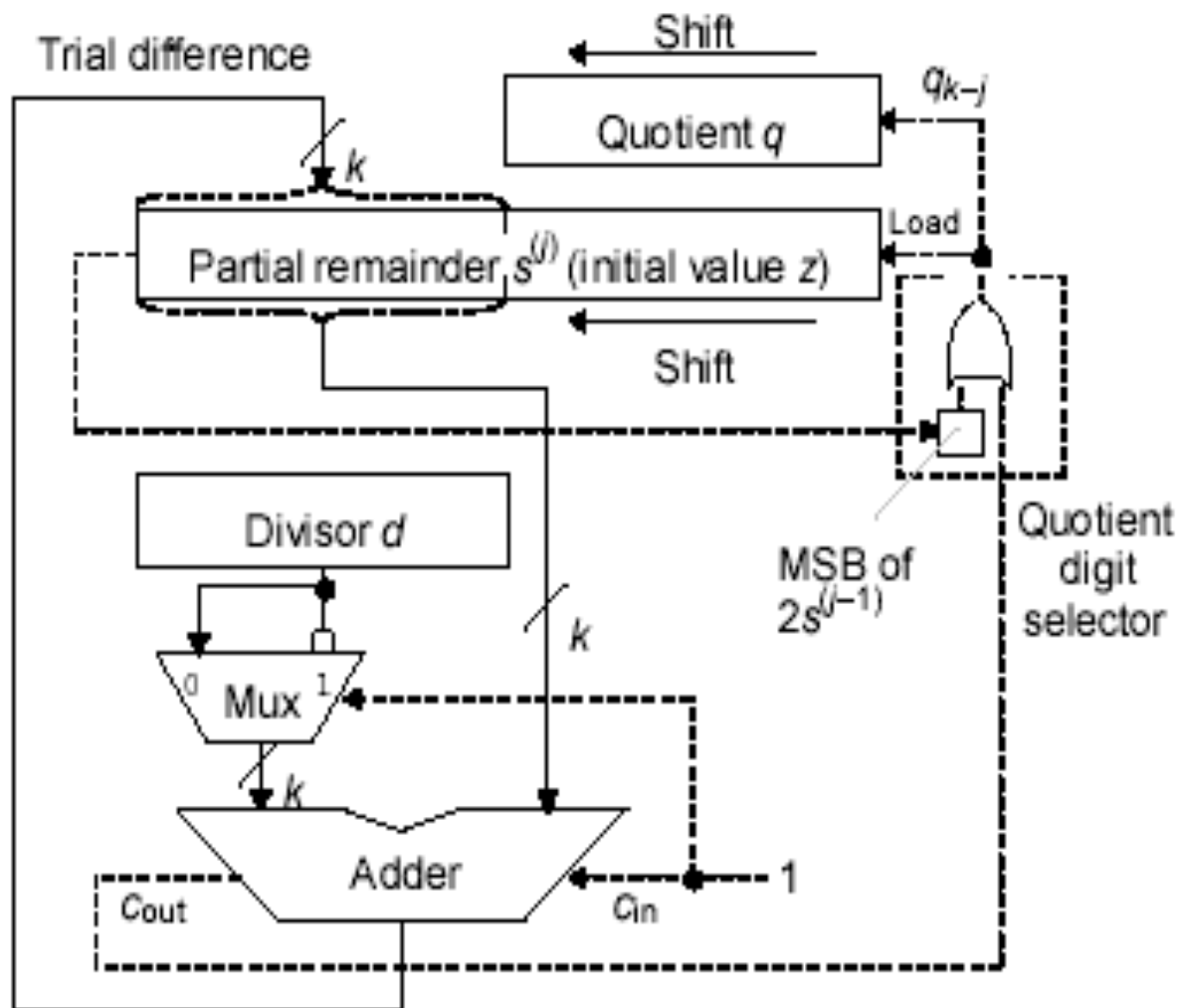
Restoring Unsigned Integer Division

```
s(0) = z
for j = 1 to k
  if 2 s(j-1) - 2k d > 0
    qk-j = 1
    s(j) = 2 s(j-1) - 2k d
  else
    qk-j = 0
    s(j) = 2 s(j-1)
  end for
```

Fig. 13.6 Example of restoring unsigned division.

z		0	1	1	1	0 1 0 1	No overflow, since: $(0111)_{\text{two}} < (1010)_{\text{two}}$
2^4d		0	1	0	1	0	
-2^4d		1	0	1	1	0	
$s^{(0)}$		0	0	1	1	1 0 1 0 1	
$2s^{(0)}$		0	1	1	1	0 1 0 1	
$+(-2^4d)$		1	0	1	1	0	
$s^{(1)}$		0	0	1	0	0 1 0 1	Positive, so set $q_3 = 1$
$2s^{(1)}$		0	1	0	0	1 0 1	
$+(-2^4d)$		1	0	1	1	0	
$s^{(2)}$		1	1	1	1	1 0 1	Negative, so set $q_2 = 0$ and restore
$s^{(2)}=2s^{(1)}$		0	1	0	0	1 0 1	
$2s^{(2)}$		1	0	0	1	0 1	
$+(-2^4d)$		1	0	1	1	0	
$s^{(3)}$		0	1	0	0	0 1	Positive, so set $q_1 = 1$
$2s^{(3)}$		1	0	0	0	1	
$+(-2^4d)$		1	0	1	1	0	
$s^{(4)}$		0	0	1	1	1	Positive, so set $q_0 = 1$
s						0 1 1 1	
q						1 0 1 1	

Fig. 13.5 Shift/subtract sequential restoring divider.



Non-Restoring Unsigned Integer Division

Non-Restoring Unsigned Integer Division

$$s^{(0)} = z$$

$$s^{(1)} = 2 s^{(0)} - 2^k d$$

for $j = 2$ to k

if $s^{(j-1)} \geq 0$

$$q_{k-(j-1)} = 1$$

$$s^{(j)} = 2 s^{(j-1)} - 2^k d$$

else

$$q_{k-(j-1)} = 0$$

$$s^{(j)} = 2 s^{(j-1)} + 2^k d$$

end for

if $s^{(k)} \geq 0$

$$q_0 = 1$$

else

$$q_0 = 0$$

Correction step

Non-Restoring Unsigned Integer Division

Correction step

$$s = 2^{-k} \cdot s^{(k)}$$

$$z = q d + s$$

$$z, q, d \geq 0 \quad s < 0$$

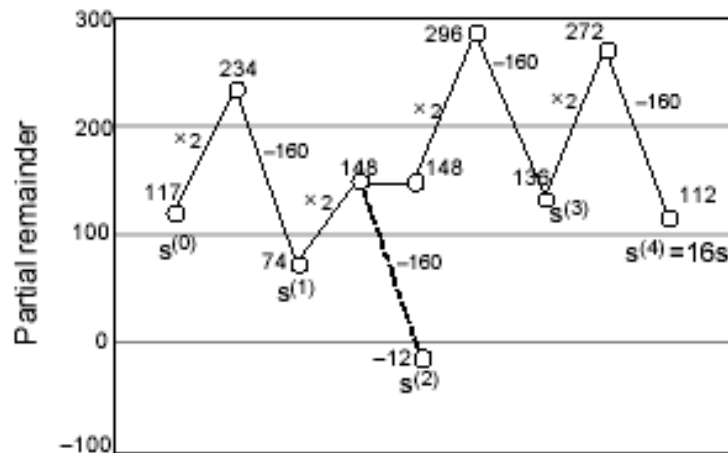
$$\begin{aligned} z &= (q-1) d + (s+d) \\ z &= q' d + s' \end{aligned}$$

Example of nonrestoring unsigned division

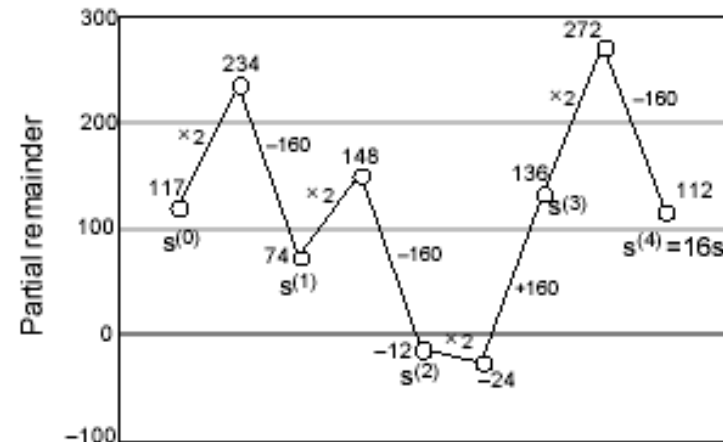
z	0	1	1	1	0	1	0	1	No overflow, since:	
2^4d	0	1	0	1	0					$(0111)_{\text{two}} < (1010)_{\text{two}}$
-2^4d	1	0	1	1	0					
$s^{(0)}$	0	0	1	1	1	0	1	0	1	
$2s^{(0)}$	0	1	1	1	0	1	0	1	Positive,	
$+(-2^4d)$	1	0	1	1	0					so subtract
$s^{(1)}$	0	0	1	0	0	1	0	1		
$2s^{(1)}$	0	1	0	0	1	0	1	Positive, so set $q_3 = 1$		
$+(-2^4d)$	1	0	1	1	0					and subtract
$s^{(2)}$	1	1	1	1	1	0	1			
$2s^{(2)}$	1	1	1	1	0	1				
$+2^4d$	0	1	0	1	0					Negative, so set $q_2 = 0$
$s^{(3)}$	0	1	0	0	0	1				
$2s^{(3)}$	1	0	0	0	1					
$+(-2^4d)$	1	0	1	1	0					Positive, so set $q_1 = 1$
$s^{(4)}$	0	0	1	1	1					Positive, so set $q_0 = 1$
s						0	1	1	1	
q						1	0	1	1	

Partial remainder variations for restoring and nonrestoring division

Example
 $(0\ 1\ 1\ 1\ 0\ 1)_{two} / (1\ 0\ 1\ 0)_{two}$
 $(117)_{ten} / (10)_{ten}$



(a) Restoring



(b) Nonrestoring

Non-Restoring Unsigned Integer Division

Justification

$$s^{(j-1)} \geq 0$$

$$2 s^{(j-1)} - 2^k d < 0$$

$$2 (2 s^{(j-1)}) - 2^k d \geq 0$$

Restoring division

$$s^{(j)} = 2 s^{(j-1)}$$

$$\begin{aligned} s^{(j+1)} &= 2 s^{(j)} - 2^k d = \\ &= 4 s^{(j-1)} - 2^k d \end{aligned}$$

Non-Restoring division

$$s^{(j)} = 2 s^{(j-1)} - 2^k d$$

$$\begin{aligned} s^{(j+1)} &= 2 s^{(j)} + 2^k d = \\ &= 2 (2 s^{(j-1)} - 2^k d) + 2^k d = \\ &= 4 s^{(j-1)} - 2^k d \end{aligned}$$

Convergence of the Partial Quotient to q

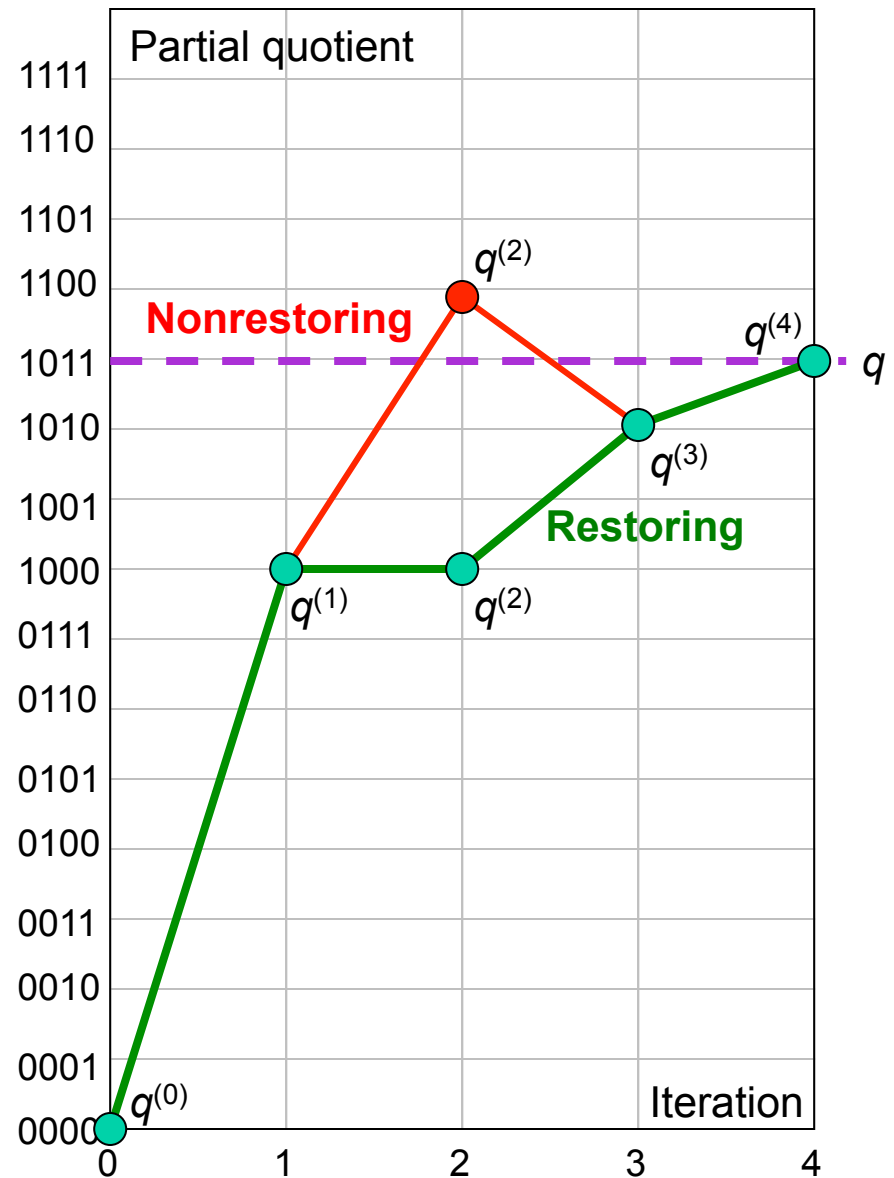
Example

$$(0\ 1\ 1\ 1\ 0\ 1\ 0\ 1)_{\text{two}} / (1\ 0\ 1\ 0)_{\text{two}}$$

$$(117)_{\text{ten}} / (10)_{\text{ten}} = (11)_{\text{ten}} = (1011)_{\text{two}}$$

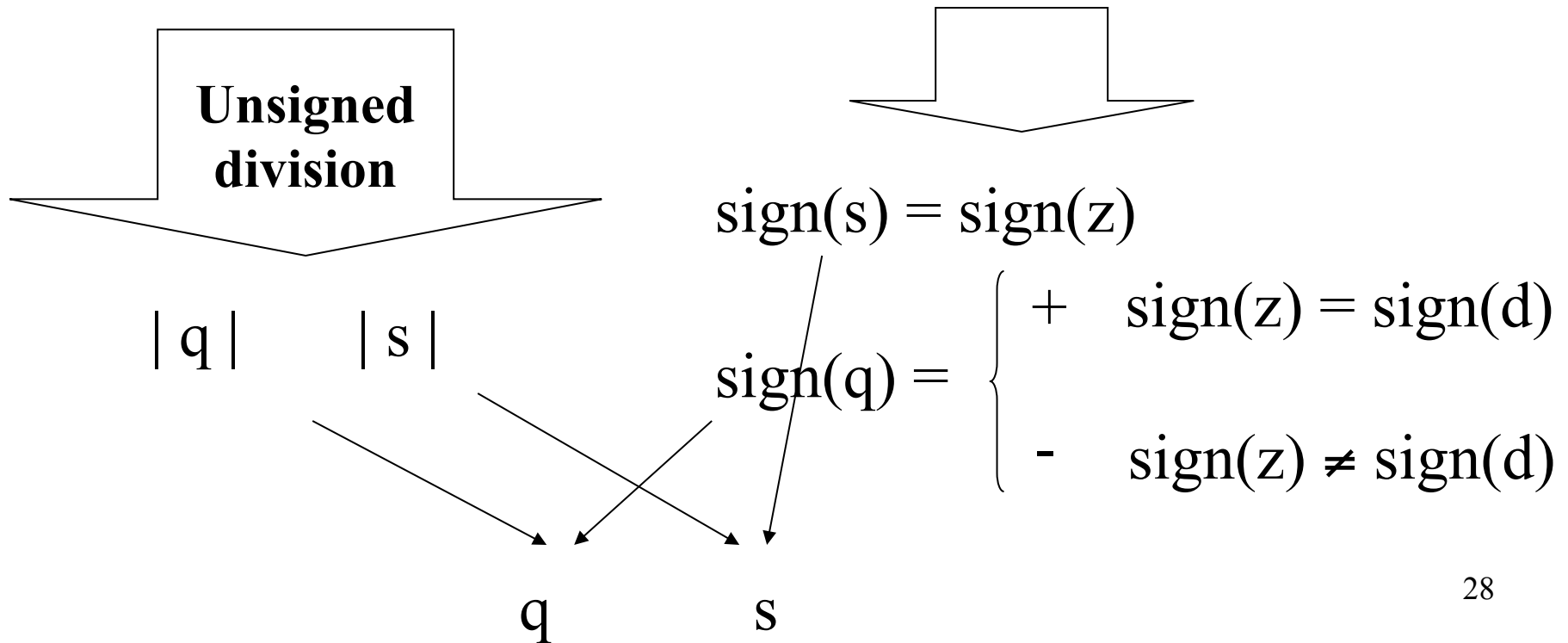
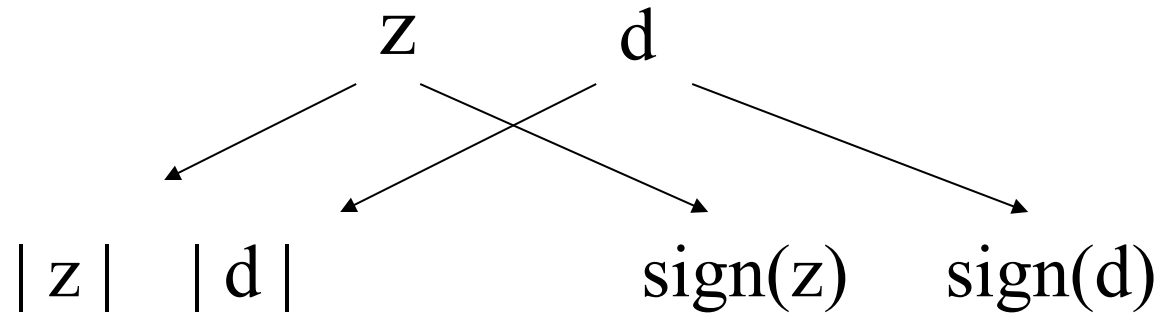
In restoring division, the partial quotient converges to q from below

In nonrestoring division, the partial quotient may overshoot q , but converges to it after some oscillations



Signed Integer Division

Signed Integer Division



Examples of Signed Integer Division

Examples of division with signed operands

$$z = 5 \quad d = 3 \quad \Rightarrow \quad q = 1 \quad s = 2$$

$$z = 5 \quad d = -3 \quad \Rightarrow \quad q = -1 \quad s = 2$$

$$z = -5 \quad d = 3 \quad \Rightarrow \quad q = -1 \quad s = -2$$

$$z = -5 \quad d = -3 \quad \Rightarrow \quad q = 1 \quad s = -2$$

Magnitudes of q and s are unaffected by input signs
Signs of q and s are derivable from signs of z and d

Non-Restoring Signed Integer Division

Non-Restoring Signed Integer Division

$$s^{(0)} = z$$

for $j = 1$ to k

if $\text{sign}(s^{(j-1)}) == \text{sign}(d)$

$$q_{k-j} = 1$$

$$s^{(j)} = 2 s^{(j-1)} - 2^k d = 2 s^{(j-1)} - q_{k-j} (2^k d)$$

else

$$q_{k-j} = -1$$

$$s^{(j)} = 2 s^{(j-1)} + 2^k d = 2 s^{(j-1)} - q_{k-j} (2^k d)$$

end for

$q = \text{BSD_2's_comp_conversion}(q)$

Correction_step

Non-Restoring Signed Integer Division

Correction step

$$s = 2^{-k} \cdot s^{(k)}$$

$$z = q d + s$$

$$\text{sign}(s) = \text{sign}(z)$$

$$\begin{aligned} z &= (q-1) d + (s+d) \\ z &= q' d + s' \end{aligned}$$

$$\begin{aligned} z &= (q+1) d + (s-d) \\ z &= q'' d + s'' \end{aligned}$$

Example of nonrestoring signed division

z	0 0 1 0 0 0 0 1
2^4d	1 1 0 0 1
-2^4d	0 0 1 1 1
$s^{(0)}$	0 0 0 1 0 0 0 0 1
$2s^{(0)}$	0 0 1 0 0 0 0 1
$+2^4d$	1 1 0 0 1
$s^{(1)}$	1 1 1 0 1 0 0 1
$2s^{(1)}$	1 1 0 1 0 0 1
$+(-2^4d)$	0 0 1 1 1
$s^{(2)}$	0 0 0 0 1 0 1
$2s^{(2)}$	0 0 0 1 0 1
$+2^4d$	1 1 0 0 1
$s^{(3)}$	1 1 0 1 1 1
$2s^{(3)}$	1 0 1 1 1
$+(-2^4d)$	0 0 1 1 1
$s^{(4)}$	1 1 1 1 0
$+(-2^4d)$	0 0 1 1 1
$s^{(4)}$	0 0 1 0 1
s	0 1 0 1
q	-1 1 -1 1

$\text{sign}(s^{(0)}) \neq \text{sign}(d)$,
so set $q_3 = -1$ and add

$\text{sign}(s^{(1)}) = \text{sign}(d)$,
so set $q_2 = 1$ and subtract

$\text{sign}(s^{(2)}) \neq \text{sign}(d)$,
so set $q_1 = -1$ and add

$\text{sign}(s^{(3)}) = \text{sign}(d)$,
so set $q_0 = 1$ and subtract

$\text{sign}(s^{(4)}) \neq \text{sign}(z)$,
so perform corrective subtraction

$p =$ 0 1 0 1 Shift, compl MSB
 1 1 0 1 1 Add 1 to correct
 1 1 0 0 Check: $33/(-7) = -4$

BSD \rightarrow 2's Complement Conversion

$$\begin{aligned}
 q &= (q_{k-1} q_{k-2} \cdots q_1 q_0)_{\text{BSD}} = \\
 &= \overline{(p_{k-1} p_{k-2} \cdots p_1 p_0 1)}_{2' \text{ s complement}}
 \end{aligned}$$

where

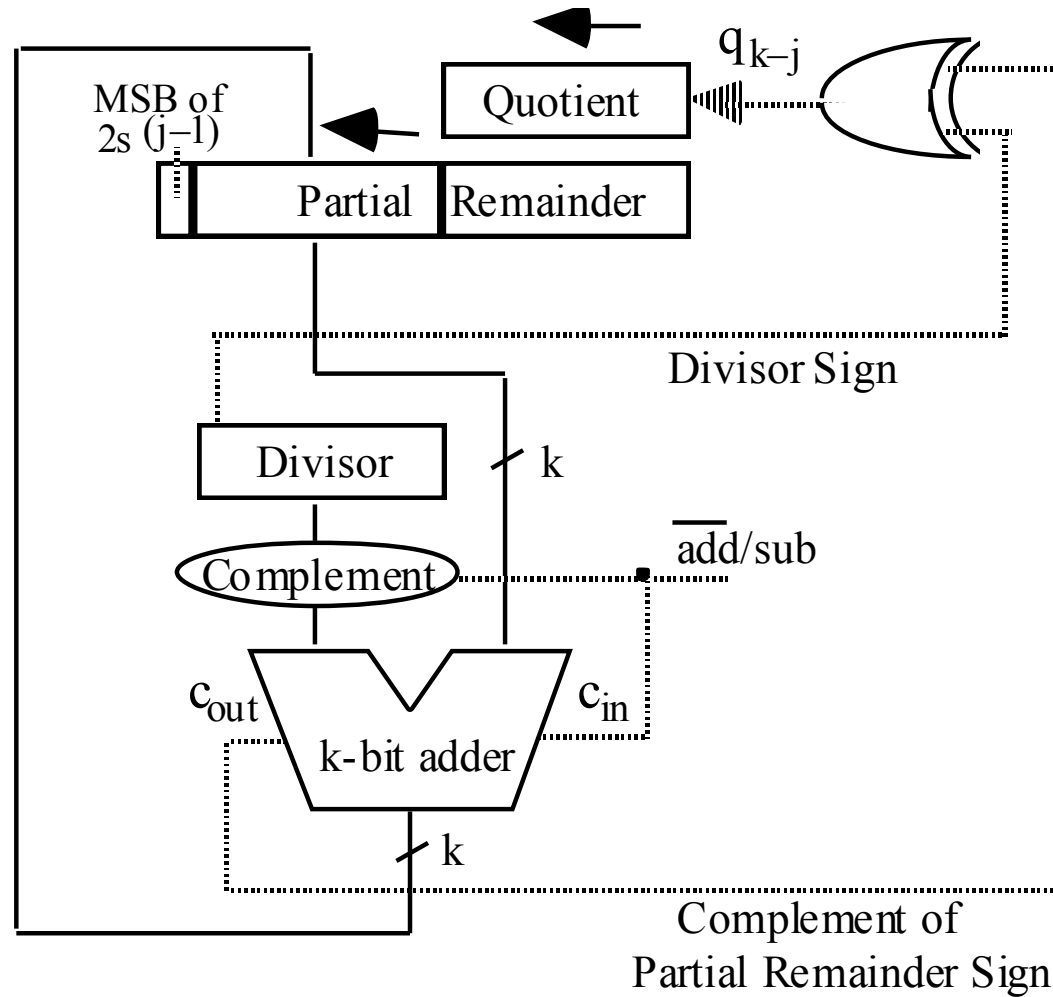
q_i	p_i
-1	0
1	1

Example:

$$\begin{array}{rcl}
 q_{\text{BSD}} & & 1 \ -1 \ 1 \ 1 \\
 p & & 1 \ 0 \ 1 \ 1 \\
 q_{2' \text{ s comp}} & & \underline{0 \ 0 \ 1 \ 1 \ 1} = 0 \ 1 \ 1 \ 1
 \end{array}$$

no overflow if $p_{k-2} = \overline{p_{k-1}}$ ($q_{k-1} \neq q_{k-2}$)

Nonrestoring Hardware Divider



Multiply/Divide Unit

Multiply-Divide Unit

The control unit proceeds through necessary steps for multiplication or division (including using the appropriate shift direction)

The slight speed penalty owing to a more complex control unit is insignificant

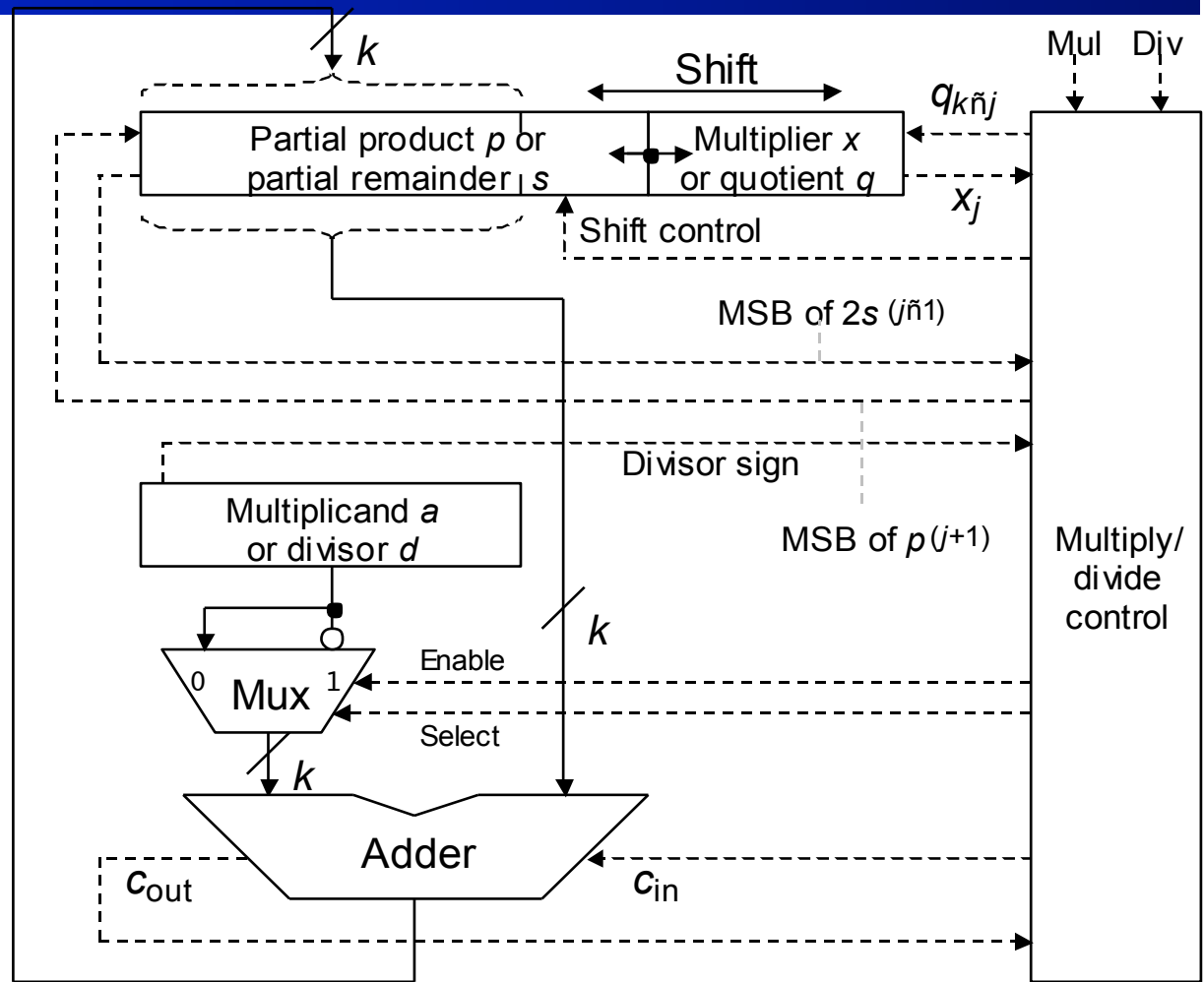


Fig. 15.9 Sequential radix-2 multiply/divide unit.